

A Perfect Path from Computational Biology to Quantum Computing

Celina Miraglia Herrera de Figueiredo





Celina, 1991



Luerbio, 1998



Simone, 2002



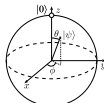
Vânia, 2004



Cláudia, 2005



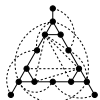
Vinicius, 2006



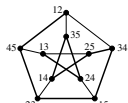
Luis Kowada, 2006



Rodrigo, 2007



Rafael Bernardo, 2008



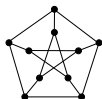
Leticia, 2009



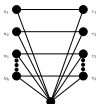
Raphael Machado, 2010



André, 2013



Diana, 2013



Hélio, 2014



Luis Felipe, 2017



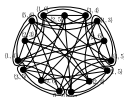
Ana Luísa, 2017



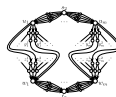
Alexandre, 2020



Edineço, 2021

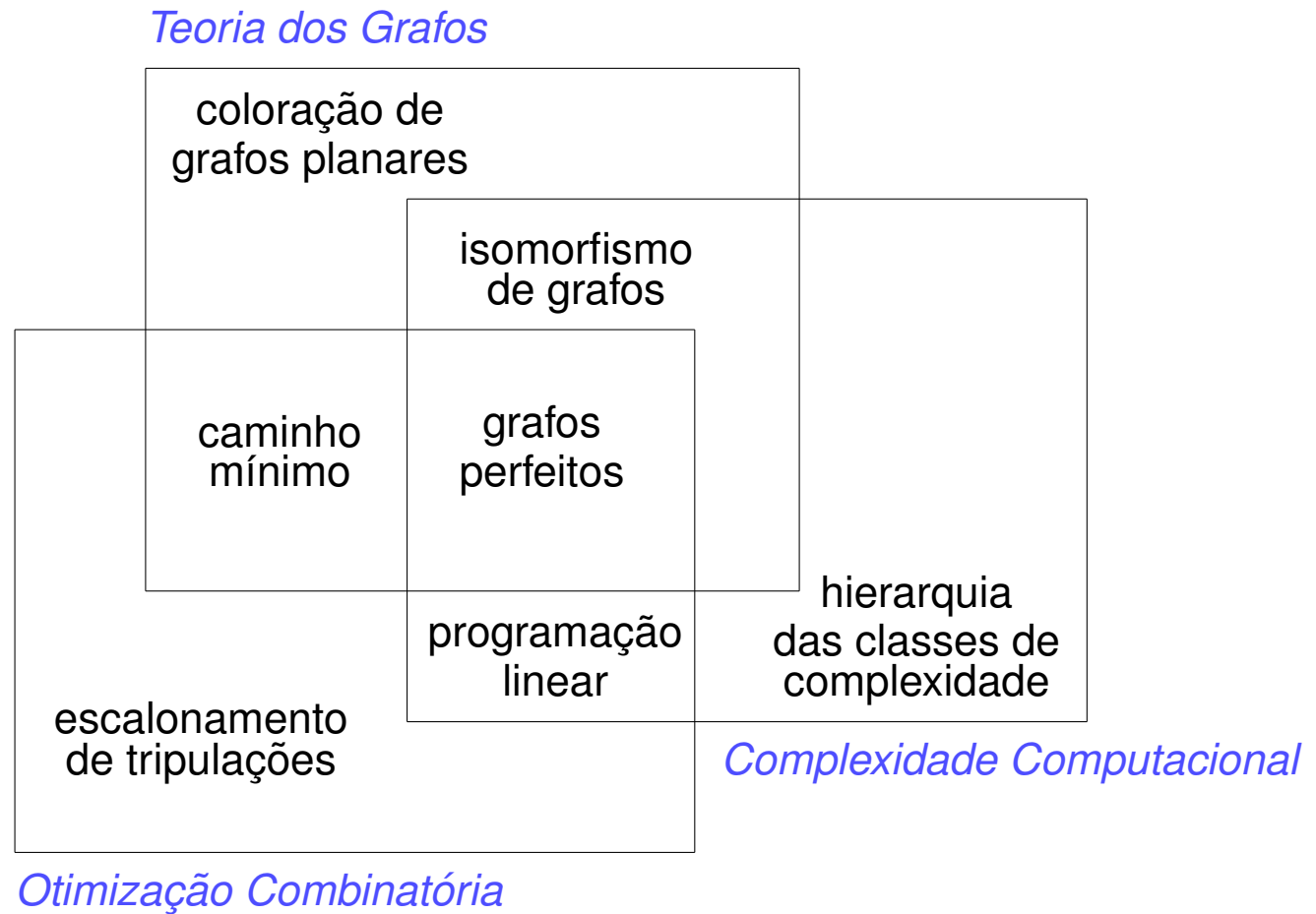


Caroline, 2021

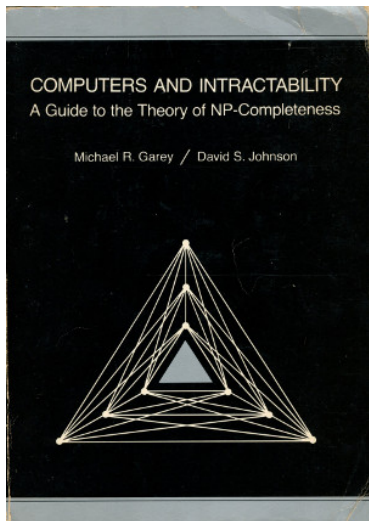


Alexander, 2022

Origem e desenvolvimento da área de pesquisa



The Guide – Computers and Intractability



“Despite that 23 years have passed since its publication, I consider Garey and Johnson the single most important book on my office bookshelf. Every computer scientist should have this book on their shelves as well. NP-completeness is the single most important concept to come out of theoretical computer science and no book covers it as well as Garey and Johnson.”

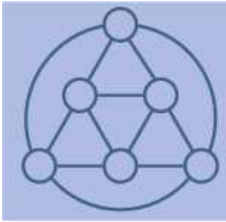
Lance Fortnow, “Great Books: Computers and Intractability: A Guide to the Theory of NP-Completeness”

Ongoing Guide – Graph Restrictions and Their Effect

GRAPH CLASS	MEMBER	INDSET	CLIQUE	CLIPAR	CHRNUM	CHRIND	HAMCIR	DOMSET	MAXCUT	STTREE	GRAISO
Trees/Forests	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	P [GJ]	P [T]	P [GJ]
Almost Trees (k)	P	P [24]	P [T]	P?	P?	P?	P?	P [45]	P?	P?	P?
Partial k -Trees	P [2]	P [1]	P [T]	P?	P [1]	O?	P [3]	P [3]	P?	P?	O?
Bandwidth- k	P [68]	P [64]	P [T]	P?	P [64]	P?	P?	P [64]	P [64]	P?	P [58]
Degree- k	P [T]	N [GJ]	P [T]	N [GJ]	N [GJ]	N [49]	N [GJ]	N [GJ]	N [GJ]	N [GJ]	P [58]
Planar	P [GJ]	N [GJ]	P [T]	N [10]	N [GJ]	O	N [GJ]	N [GJ]	P [GJ]	N [35]	P [GJ]
Series Parallel	P [79]	P [75]	P [T]	P?	P [74]	P [74]	P [74]	P [54]	P [GJ]	P [82]	P [GJ]
Outerplanar	P	P [6]	P [T]	P [6]	P [67]	P [67]	P [T]	P [6]	P [GJ]	P [81]	P [GJ]
Halin	P	P [6]	P [T]	P [6]	P [74]	P [74]	P [T]	P [6]	P [GJ]	P?	P [GJ]
k -Outerplanar	P	P [6]	P [T]	P [6]	P [6]	O?	P [6]	P [6]	P [GJ]	P?	P [GJ]
Grid	P	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	N [51]	N [55]	P [T]	N [35]	P [GJ]
$K_{3,3}$ -Free	P [4]	N [GJ]	P [T]	N [10]	N [GJ]	O?	N [GJ]	N [GJ]	P [5]	N [GJ]	O?
Thickness- k	N [60]	N [GJ]	P [T]	N [10]	N [GJ]	N [49]	N [GJ]	N [GJ]	N [7]	N [GJ]	O?
Genus- k	P [34]	N [GJ]	P [T]	N [10]	N [GJ]	O?	N [GJ]	N [GJ]	O?	N [GJ]	P [61]
Perfect	O!	P [42]	P [42]	P [42]	P [42]	O?	N [1]	N [14]	O?	N [GJ]	I [GJ]
Chordal	P [76]	P [40]	P [40]	P [40]	P [40]	O?	N [22]	N [14]	O?	N [83]	I [GJ]
Split	P [40]	P [40]	P [40]	P [40]	P [40]	O?	N [22]	N [19]	O?	N [83]	I [15]
Strongly Chordal	P [31]	P [40]	P [40]	P [40]	P [40]	O?	O?	P [32]	O?	P [83]	O?
Comparability	P [40]	P [40]	P [40]	P [40]	P [40]	O?	N [1]	N [28]	O?	N [GJ]	I [GJ]
Bipartite	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	N [1]	N [28]	P [T]	N [GJ]	I [GJ]
Permutation	P [40]	P [40]	P [40]	P [40]	P [40]	O?	O	P [33]	O?	P [23]	P [21]
Cographs	P [T]	P [40]	P [40]	P [40]	P [40]	O?	P [25]	P [33]	O?	P [23]	P [25]
Undirected Path	P [39]	P [40]	P [40]	P [40]	P [40]	O?	O?	N [16]	O?	O?	I [GJ]
Directed Path	P [38]	P [40]	P [40]	P [40]	P [40]	O?	O?	P [16]	O?	P [83]	O?
Interval	P [17]	P [44]	P [44]	P [44]	P [44]	O?	P [53]	P [16]	O?	P [83]	P [57]
Circular Arc	P [78]	P [44]	P [50]	P [44]	N [36]	O?	O?	P [13]	O?	P [83]	O?
Circle	P [71]	P [GJ]	P [50]	O?	N [36]	O?	P [12]	O?	O?	P [70]	O?
Proper Circ. Arc	P [77]	P [44]	P [50]	P [44]	P [66]	O?	P [12]	P [13]	O?	P [83]	O?
Edge (or Line)	P [47]	P [GJ]	P [T]	N [GJ]	N [49]	O?	N [11]	N [GJ]	O?	N [70]	I [15]
Claw-Free	P [T]	P [63]	O?	N [GJ]	N [49]	O?	N [11]	N [GJ]	O?	N [70]	I [15]

The updated NP-Completeness Column: An Ongoing Guide table 35 years later

GRAPH CLASS	MEMBER	INDSET	CLIQUE	CLIPAR	CHRNUM	CHRIND	HAMCIR	DOMSET	MAXCUT	STTREE	GRAPHISO
TREES/FORESTS	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	P [GJ]	P [T]	P [GJ]
ALMOST TREES (K)	P [OG]	P [OG]	P [T]	P [105]	P [5]	P [17]	P [5]	P [5]	P [20]	P [76]	P [17]
PARTIAL K-TREES	P [OG]	P [5]	P [T]	P [105]	P [5]	P [17]	P [5]	P [5]	P [20]	P [76]	P [17]
BANDWIDTH-K	P [OG]	P [OG]	P [T]	P [105]	P [5]	P [17]	P [5]	P [5]	P [OG]	P [76]	P [OG]
DEGREE-K	P [T]	N [GJ]	P [T]	N [29]	N [GJ]	N [OG]	N [GJ]	N [GJ]	N [GJ]	N [GJ]	P [OG]
PLANAR	P [GJ]	N [GJ]	P [T]	N [78]	N [GJ]	O	N [GJ]	N [GJ]	P [GJ]	N [OG]	P [GJ]
SERIES PARALLEL	P [OG]	P [OG]	P [T]	P [105]	P [5]	P [17]	P [5]	P [OG]	P [GJ]	P [OG]	P [GJ]
OUTERPLANAR	P [OG]	P [OG]	P [T]	P [OG]	P [OG]	P [OG]	P [T]	P [OG]	P [GJ]	P [OG]	P [GJ]
HALIN	P [OG]	P [OG]	P [T]	P [OG]	P [5]	P [17]	P [T]	P [OG]	P [GJ]	P [118]	P [GJ]
K-OUTERPLANAR	P [OG]	P [OG]	P [T]	P [OG]	P [5]	P [17]	P [OG]	P [OG]	P [GJ]	P [76]	P [GJ]
GRID	P [OG]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	N [OG]	N [32]	P [T]	N [OG]	P [GJ]
K _{3,3} -FREE*	P [OG]	N [GJ]	P [T]	N [78]	N [GJ]	O?	N [GJ]	N [GJ]	P [OG]	N [GJ]	P [40]
THICKNESS-K	N [OG]	N [GJ]	P [T]	N [78]	N [GJ]	N [OG]	N [GJ]	N [GJ]	N [119]	N [GJ]	I [RJ]
GENUS-K	P [OG]	N [GJ]	P [T]	N [78]	N [GJ]	O?	N [GJ]	N [GJ]	O?	N [GJ]	P [OG]
PERFECT	P [34]	P [OG]	P [OG]	P [OG]	P [OG]	N [28]	N [OG]	N [OG]	N [20]	N [GJ]	I [84]
CHORDAL	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	N [93]	N [OG]	N [20]	N [OG]	I [84]
SPLIT	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	N [93]	N [OG]	N [20]	N [OG]	I [108]
STRONGLY CHORDAL	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	N [93]	P [OG]	N [109]	P [OG]	I [111]
COMPARABILITY	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	N [28]	N [OG]	N [94]	N [102]	N [GJ]	I [22]
BIPARTITE	P [T]	P [GJ]	P [T]	P [GJ]	P [T]	P [GJ]	N [OG]	N [94]	P [T]	N [GJ]	I [22]
PERMUTATION	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	P [44]	P [OG]	N [120]	P [OG]	P [OG]
COGRAPHS	P [T]	P [OG]	P [OG]	P [OG]	P [OG]	O?	P [OG]	P [OG]	P [20]	P [OG]	P [OG]
UNDIRECTED Path	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	N [13]	N [OG]	N [20]	N [RJ]	I [22]
DIRECTED PATH	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	N [99]	P [OG]	N [11]	P [OG]	P [7]
INTERVAL	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	P [OG]	P [OG]	N [11]	P [OG]	P [OG]
CIRCULAR ARC	P [OG]	P [OG]	P [OG]	P [OG]	N [OG]	O?	P [106]	P [OG]	N [11]	P [11]	P [80]
CIRCLE	P [OG]	P [GJ]	P [OG]	N [73]	N [OG]	O?	N [39]	N [71]	N [26]	P [OG]	P [68]
PROPER CIRC. ARC	P [OG]	P [OG]	P [OG]	P [OG]	P [OG]	O?	P [OG]	P [OG]	O?	P [11]	P [82]
EDGE (OR LINE)	P [OG]	P [GJ]	P [T]	N [95]	N [OG]	N [28]	N [OG]	N [GJ]	P [59]	N [19]	I [OG]
CLAW-FREE	P [T]	P [OG]	N [103]	N [85]	N [OG]	N [28]	N [OG]	N [GJ]	N [20]	N [19]	I [OG]



LAGOS 2009

V Latin-American Algorithms, Graphs and Optimization Symposium



The P vs. NP-complete dichotomy of some challenging problems in graph theory

Celina de Figueiredo

Universidade Federal do Rio de Janeiro

Brazil

November 2009

Two long-standing problems in graph theory

Perfect graphs: Chvátal's SKEW PARTITION is polynomial

Intersection graphs: Roberts–Spencer's CLIQUE GRAPH is NP-complete

Both SKEW PARTITION and CLIQUE GRAPH proved to be in NP when their classification into P or NP-complete was proposed

V. Chvátal – *J. Combin. Theory Ser. B* 1985

F. Roberts, J. Spencer – *J. Combin. Theory Ser. B* 1971

The three nonempty part problem

Full dichotomy for the RECOGNITION PROBLEM:

STABLE CUTSET, 3-COLORING are the only NP-complete

T. Feder, P. Hell, S. Klein, R. Motwani – *SIAM J. Discrete Math.* 2003

Full dichotomy for the SANDWICH PROBLEM:

61 interesting problems: 19 NP-complete, 42 polynomial

HOMOGENEOUS SET SANDWICH PROBLEM is polynomial

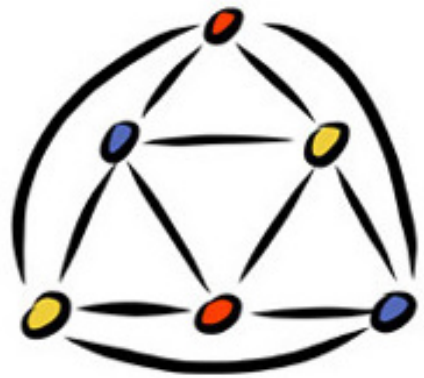
CLIQUE CUTSET SANDWICH PROBLEM is NP-complete

Full dichotomy for the GENERALIZED SPLIT GRAPH SANDWICH PROBLEM:

(2,1)-GRAPH SANDWICH PROBLEM is NP-complete

“The polynomial dichotomy for three nonempty part sandwich problems”

Discrete Appl. Math. 2009 (with Rafael Teixeira, Simone Dantas)



LAGOS 2017

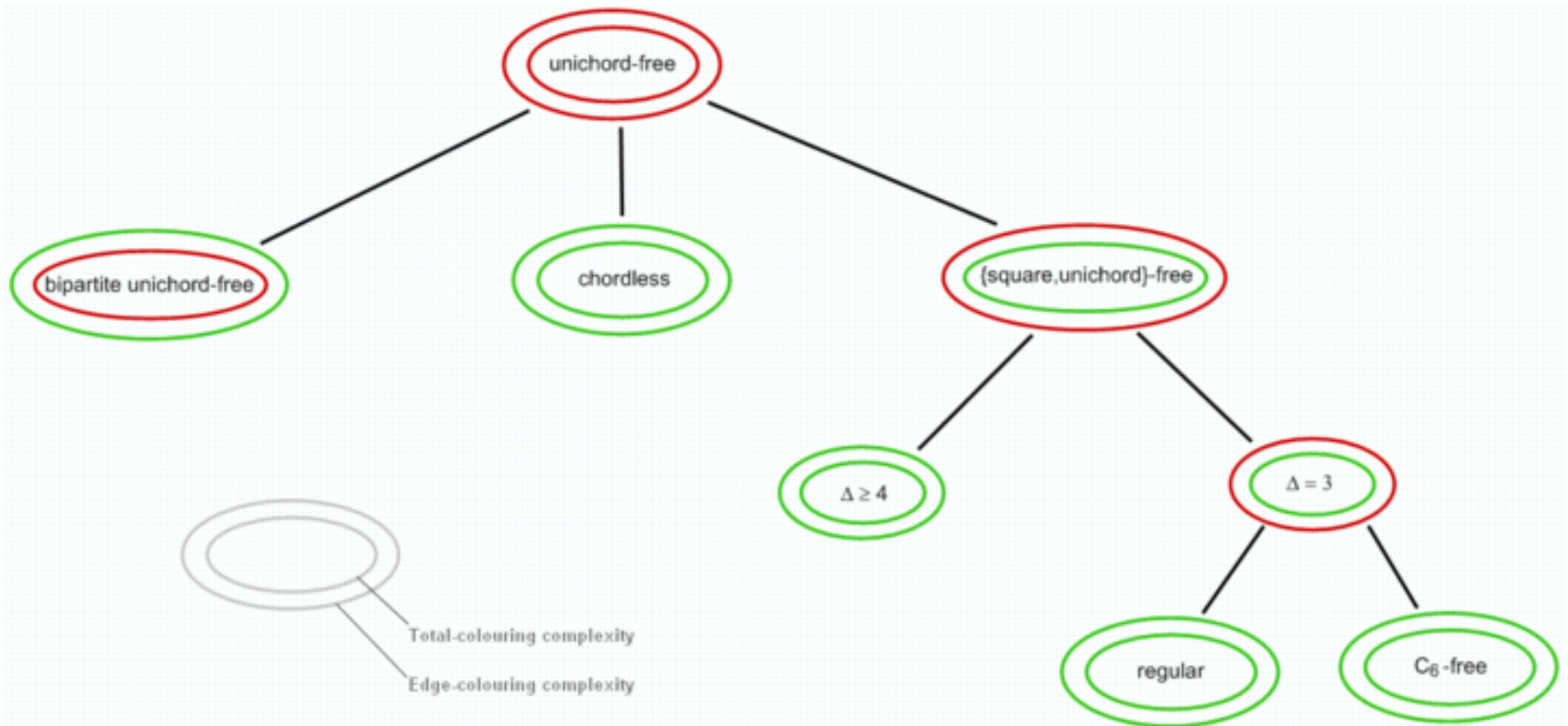
Complexity-separating graph classes for vertex, edge and total coloring

Celina de Figueiredo



COPPE
UFRJ

Edge and total coloring complexity-separating classes



When restricted to $\{\text{square, unichord}\}$ -free graphs, edge coloring is **NP-complete** whereas total coloring is **polynomial**

Complexity restricted to unichord-free and special subclasses

Colouring problem \ class	General	Unichord-free	$\{\square, \text{unichord}\}$ -free	$\{\Delta, \text{unichord}\}$ -free
Vertex-col.	\mathcal{NPC} [14]	\mathcal{P} [26]	\mathcal{P} [26]	\mathcal{P} [26]
Edge-col.	\mathcal{NPC} [13]	\mathcal{NPC} [18]	\mathcal{NPC} [18]	\mathcal{NPC} [18]
Total-col.	\mathcal{NPC} [21]	\mathcal{NPC} [17]	\mathcal{P} [16,17]	\mathcal{NPC} [17]
Clique-col.	$\Sigma_2^P \mathcal{C}$ [20]	\mathcal{P}	\mathcal{P}	$\mathcal{P} (\kappa = \chi)$
Biclique-col.	$\Sigma_2^P \mathcal{C}$ [10]	\mathcal{P}	\mathcal{P}	$\mathcal{P} (\kappa_{\mathbf{B}} = 2)$

[10] M. Groshaus, F. Soulignac, P. Terlisky – *J. Graph Algorithms Appl.* 2014

[20] D. Marx – *Theoret. Comput. Sci.* 2011

“Efficient algorithms for clique-colouring and biclique-colouring unichord-free graphs”
Algorithmica 2017 (with Hélio Macedo and Raphael Machado)

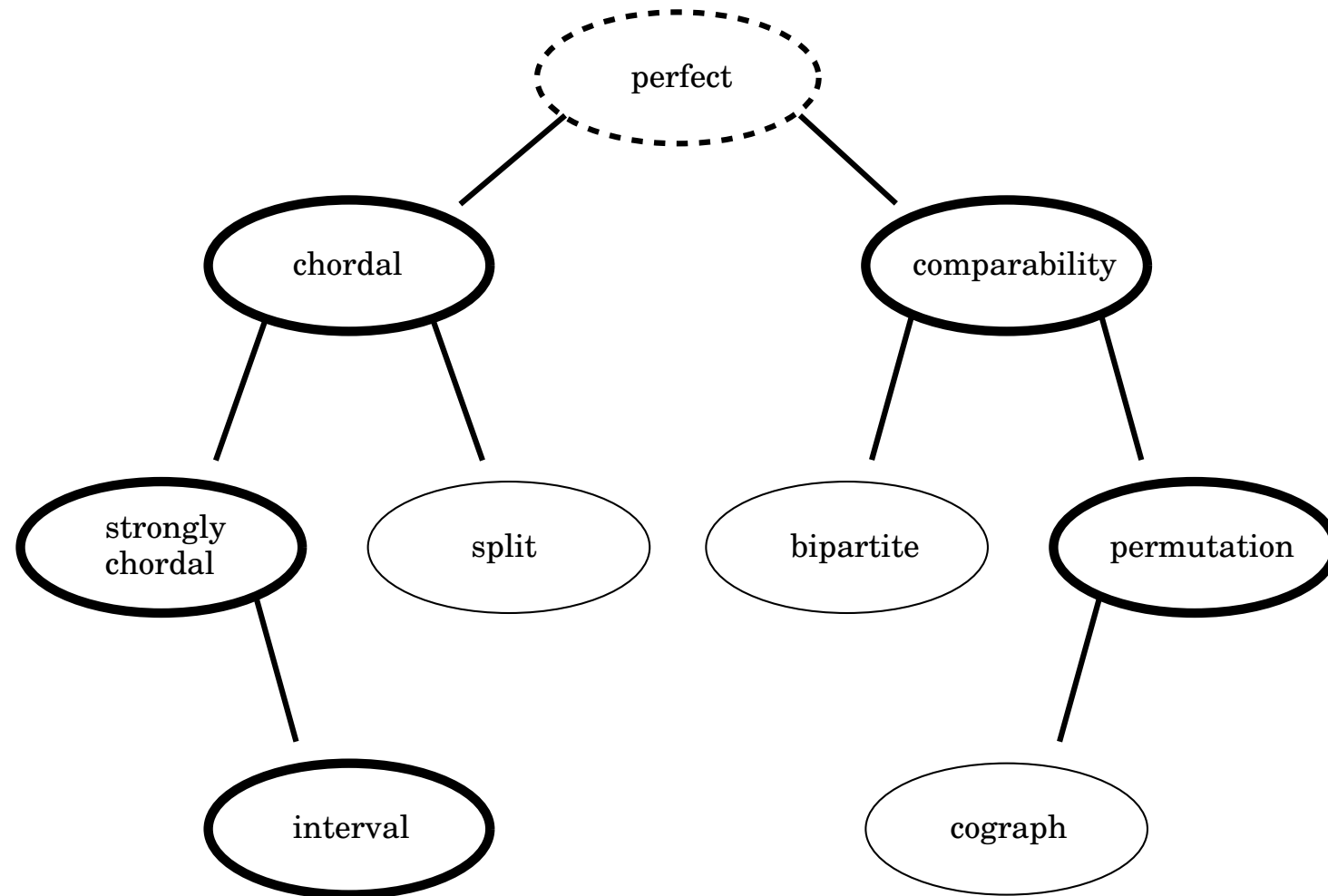
Every graph is easy or hard: dichotomy theorems for graph problems

Dániel Marx¹

¹Institute for Computer Science and Control,
Hungarian Academy of Sciences (MTA SZTAKI)
Budapest, Hungary

ICGT 2014
Grenoble, France
July 3, 2014

Sandwich problems for perfect graph classes



— NP-complete

— polynomial

- - - - open

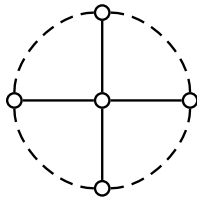
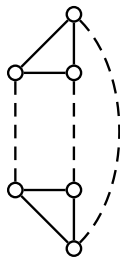
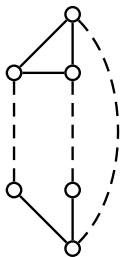
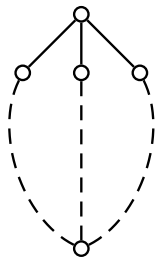
Is the not C-free easier than the C-free sandwich problem?

A trigraph (G_1, G_2) satisfies **property Π** if
there is no sandwich graph G for (G_1, G_2) which does not satisfy Π

The recognition of Berge graphs is polynomial
but
the recognition of Berge trigraphs was previously open

The imperfect graph sandwich problem is polynomial
Equivalently, recognizing Berge trigraphs is polynomial

Detecting 3-path configurations



theta and pyramid: **polynomial**

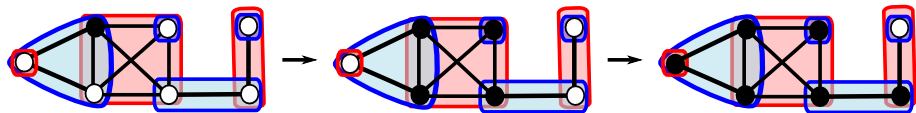
prism and wheel: **NP-complete**

The not pyramid-free sandwich problem is polynomial
but

the complexity of the pyramid-free sandwich problem is open

“The world of hereditary graph classes viewed through Truemper configurations”
by K. Vušković, in *Surveys in Combinatorics* (2013)

A quantum walker spreads across a 2-tessellation cover



The chromatic upper bound: $T(G) \leq \min\{\chi'(G), \chi(K(G))\}$

"The graph tessellation cover number: Chromatic bounds, efficient algorithms and hardness"
Theoretical Computer Science (2020) (with Alexandre Abreu, Luis Cunha, Luis Kowada,
Franklin Marquezino, Daniel Posner, Renato Portugal)

Most significant publications

- FIGUEIREDO, C. M. H. · KLEIN, S. · KOHAYAKAWA, Y. · REED, B.
Finding skew partitions efficiently
Journal of Algorithms (2000)
- FIGUEIREDO, C. M. H. · MAFFRAY, F.
Optimizing bull-free perfect graphs
SIAM Journal on Discrete Mathematics (2004)
- FARIA, L. · FIGUEIREDO, C. M. H. · SYKORA, O. · VRTO, I.
An improved upper bound on the crossing number of the hypercube
Journal of Graph Theory (2008)
- ALCON, L. · FARIA, L. · FIGUEIREDO, C. M. H. · GUTIERREZ, M.
The complexity of clique graph recognition
Theoretical Computer Science (2009)
- FIGUEIREDO, C. M. H.
The P vs. NP-complete dichotomy of some challenging problems in graph theory
Discrete Applied Mathematics (2012)

Most significant publications

- CUNHA, L. F. I. · KOWADA, L. A. B. · HAUSEN, R. A. · FIGUEIREDO, C. M. H.
A faster 1.375-approximation algorithm for sorting by transpositions
Journal of Computational Biology (2015)
- MACÊDO, H. B. · MACHADO, R. C. S. · FIGUEIREDO, C. M. H.
Hierarchical complexity of 2-clique-colouring weakly chordal graphs and perfect graphs having cliques of size at least 3
Theoretical Computer Science (2016)
- CHUDNOVSKY, M. · FIGUEIREDO, C. M. H. · SPIRKL, S.
The sandwich problem for decompositions and almost monotone properties
Algorithmica (2018)
- MELO, A. A. · FIGUEIREDO, C. M. H. · SOUZA, U. S.
A multivariate analysis of the strict terminal connection problem
Journal of Computer and System Sciences (2020)
- ABREU, A. · CUNHA, L. · FIGUEIREDO, C. · KOWADA, L. · MARQUEZINO, F. · POSNER, D. · PORTUGAL, R.
The graph tessellation cover number: Chromatic bounds, efficient algorithms and hardness
Theoretical Computer Science (2020)